TriAD: A Distributed Shared-Nothing RDF Engine based on Asynchronous Message Passing

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Motivation

RDF is ubiquitous...

Many organizations now support and publish RDF data Eg. DBpedia (400M), YAGO2 (130M), Freebase,...

How to scale and yet be efficient..?

- .. in managing and querying RDF data **Our Approach: TriAD RDF store**
- Scalability main-memory backed distributed setting
- Efficiency 1. Asynchronous query processing 2. Join-ahead pruning

Our Approach (in a nut shell)



Problems with Current Approaches

Problem 1: Synchronous processing of joins

Hadoop-based systems process joins using iterative & synchronous MapReduce jobs

Problem 2: Pruning dangling triples

- Dangling triples are the triples that appear in intermediate relations but not part of final joins
- Approaches: Graph exploration, Sideways Information Passing (SIP)
- Graph exploration works well for selective queries SIP – needs synchronization among join operators

TriAD Architecture





- Summary index (via summary bindings) is used for pruning dangling triples (join-ahead pruning)
- Asynchronous distributed query execution over RDF index

Stage 1: Join-ahead pruning (of dangling triples) Graph summarization SPARQL Query RDF Graph Cost function for SELECT ?person ?city ?prize WHERE{ **Summary Graph** optimal partitions ?person <bornIn> ?city . isΑ isA Singer Lady Gaga Barack Obama $c_{Q,n} := c_S + c_{P,n}$?city <locatedIn> USA . isA,won, bornIn locIn, bornIn bornIn/memO locIn ?person <won> ?prize. $\frac{d\left|V_{S}\right|}{\left|E_{D}\right|} \cdot c_{D} +$ Democratic Party Grammy Award bornIn Honolulu c_D SPARQL memOf C ?prize <hasName> ?name isA,won New York USA loch locln locIn memOf Texas **Optimal summary size** New Haver **Graph Exploration** locIn / governor bornIn Nobel Peace prize mem Plains $|V_S| := \sqrt{\frac{\lambda |E_D|}{d n}}$ governor 🦳 Pa George W Bush memOf ?person: P1, P2, P4 Jimmy Carter memOf,bornIn won,bornIn Republican Party ?city: P1, P2, P4 λ - tuning parameter Bindings ?prize: P2, P4 **Stage 2: Distributed & asynchronous query execution** Slave 2 Slave 1 **Global Plan** Bindings Partial Results $DHJ(R_{1,2,3,4})$ $\operatorname{DHJ}(R_{1,2,3,4})$?person 🖂 Cost:max(105,215)+130 Sharding: $R_{1,2}, R_{3,4}$

Graph Summarization & Query Processing Workflow









 $DMJ(R_{1,2})$

 $\sum^{}$ DMJ(R_{3,4})